
Introducing Design by Contract to SIDL/Babel

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Overview

- **Goals**
- **Basic Constructs**
- **Impact on SIDL/Babel**
- **Benefits for the CCA**
- **Future Work**

Why support assertions at the interface specification level?

The interface specification can provide a **simple, concise description** of the *requirements, behavior,* and *constraints*.

Generated code will **automatically ensure compliance** regardless of the underlying implementation language.

The SIDL grammar defines packages, interfaces, etc.

- Packages & Versions
- Interfaces & Classes
- Inheritance Model
- **Methods** ←
- Method Modifiers
- Intrinsic Data Types
- Parameter Modes
- And more...

*Optional
assertion
specifications
added here*

There are several types of assertions mentioned in the literature.

Type	Express...
Precondition	<ul style="list-style-type: none"> • Constraints to enable proper method function • Conditions that must be <i>true prior to</i> invocation
Postcondition	<ul style="list-style-type: none"> • Guarantees of proper method function • Conditions that must be <i>true after</i> invocation
Class Invariant	<ul style="list-style-type: none"> • Global properties of instances that must be <i>true upon</i> instance creation and preserved by all routines <i>before and after</i> every invocation
Loop Invariant	<ul style="list-style-type: none"> • Instance properties that must be <i>true prior to</i> the first execution of a loop <i>and</i> preserved by every iteration so <i>hold on</i> loop termination
Loop Variant	<ul style="list-style-type: none"> • An integer value that must be <i>non-negative prior to</i> first execution of a loop and decreased by every iteration to guarantee loop termination

Only the first two Design by Contract clauses will be added at this time.

Clause	Comment
Preconditions	<ul style="list-style-type: none"> • Specify library's requirements • Obligations on callers
Postconditions	<ul style="list-style-type: none"> • Specify library's guarantees • Obligations on callee to <i>provided</i> preconditions were satisfied <i>and no exceptions raised</i>

Method sequencing will be implemented using a sequence, or state, clause the values of which can be utilized in the pre- and post-conditions to specify method ordering.

The new clauses require a number of additions to the SIDL grammar.

- Eiffel keywords
 - Preconditions `requires`
 - Postconditions `ensures`
- Simple conditional expression operators
 - Logical `&&, ||, !`
 - Bitwise Logical `&, ^, |`
 - Relational `<, <=, ==, !=, >=, >`
 - Shift `<<, >>`
 - Additive `+, -`
 - Multiplicative `*, /, %`
- Logical grouping `()`
- Literal keywords `TRUE, FALSE, NULL, return`
- Terminals

CASC

TLD 7

The following specification snippet illustrates the use of both clauses in SIDL.

```
interface Vector {  
    Vector axpy (in Vector a, in Vector x) {  
        requires a != NULL;  
                x != NULL;  
        ensures return != NULL;  
    };  
    double norm () {  
        ensures return >= 0.0;  
    };  
};
```

Vector.sidl

Recall: If method raises an exception then *no* guarantee the ensures will be met!

CASC

TLD 8

If method sequencing were incorporated, call ordering state would be added.

```
interface Vector {  
    state { uninitialized, initialized };  
  
    void setData (in double data){  
        requires uninitialized;  
        ensures initialized;  
    };  
    ...  
}
```

First item is the initial state.

Transition to *initialized* is automatic if library call is successful and all (other) postcondition entries met.

VectorWithOrdering.sidl

Note: Sequencing constructs subject to change.

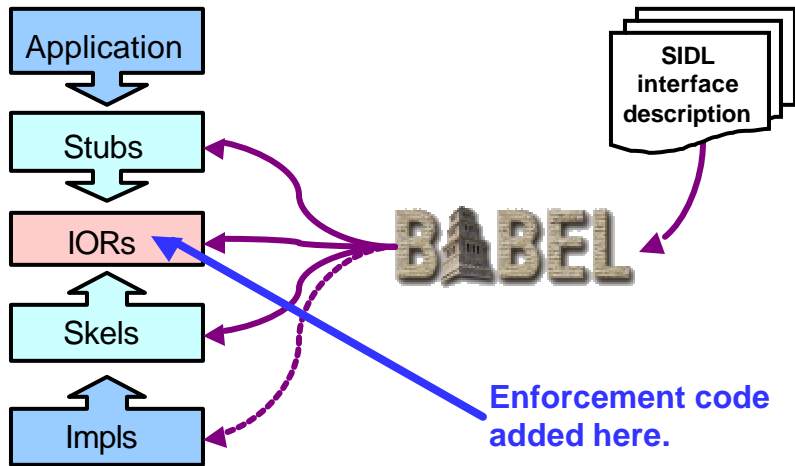
Methods that require the instance to be in a state could annotate it accordingly.

```
interface Vector {  
    ...  
    Vector axpy (in Vector a, in Vector x) {  
        requires initialized;  
        a != NULL;  
        x != NULL;  
        ensures return != NULL;  
    };  
    double norm () {  
        requires initialized;  
        ensures return >= 0.0;  
    };  
}
```

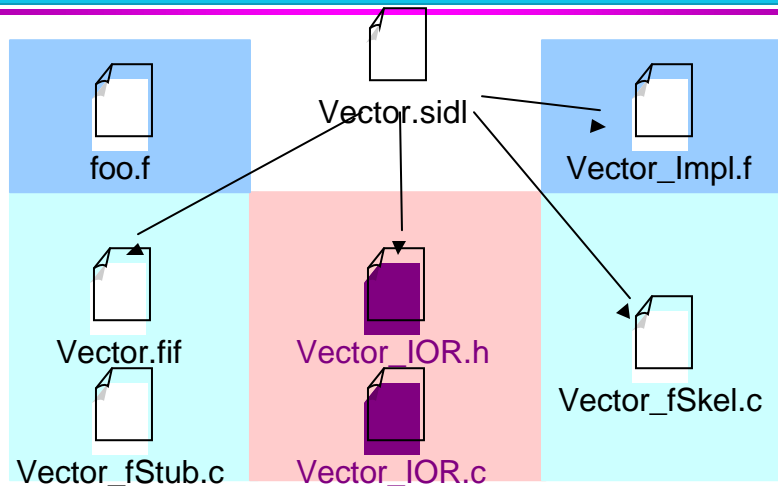
VectorWithOrdering.sidl

Note: Sequencing constructs subject to change.

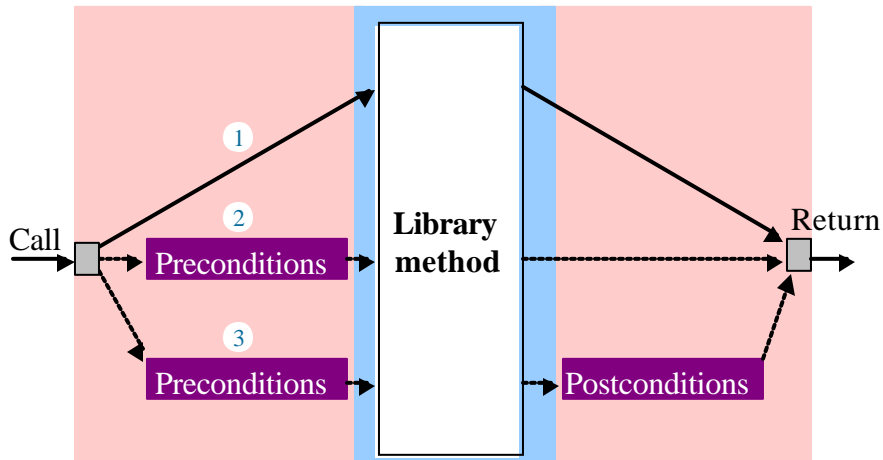
Babel takes a SIDL file and will generate expanded glue code.



The IOR files will be changed to add the generated checks.



There will be three execution paths available through the IOR.



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TLD 13

Dynamic switching between paths can be available at up to four levels.

- Instance

- `Vector.__create(/* desired setting */);`
- `Vector.__noChecks();` // Path #1
- `Vector.__checkRequires();` // Path #2
- `Vector.__checkAsserts();` // Path #3

- Class -- For all instances of a class

- Package -- For a subset of packages

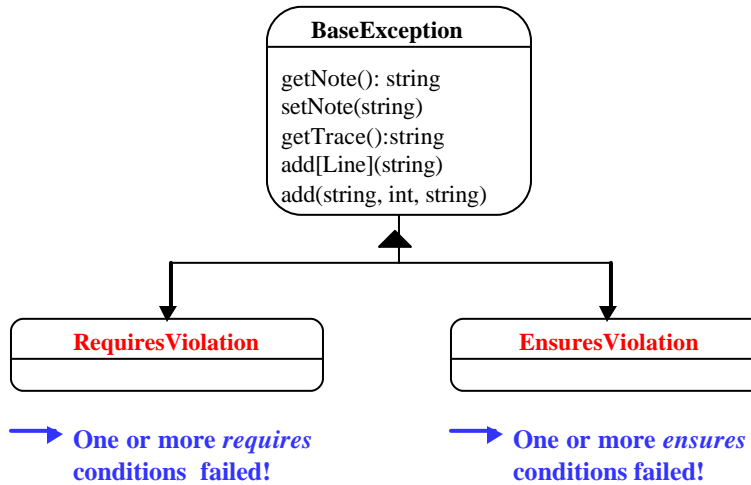
- Global -- Through the SIDL Loader

What degree of flexibility is needed?
- Dynamic switching at the Class level? Package level? Loader level?
- Regular expression support for specifying classes? Packages?

CASC

TLD 14

Violations of assertions will result in the raising of new SIDL exceptions.



Which means several parser-related files must change in the compiler.

File	Change(s)
dtds/ SIDL.dtd	Add elements for the assertion lists and conditions.
parsers/sidl/ SIDL.jj	Add support for the new grammar productions.
parsers/xml/ ParseSymbolXML.java	Add parsing for new structures from XML.
symbols/ Method.java	Add support for assertion lists.
symbols/ newclass(es).java	New file(s) associated with the assertion list productions to support the lists.

The backends must also be modified to support the IOR and stub changes.

Files	Change(s)
backend/IOR.java backend/IOR/IORHeader.java backend/IOR/IORSource.java	Add support for new built-in methods for dynamic switching and the new entry point vectors.
backend/ <i>language</i> /StubHeader.java* backend/ <i>language</i> /StubSource.java*	Add support for new built-in methods for dynamic switching.

* These or their equivalent are generally present for each supported language.

Interface-level assertions will ultimately facilitate wider reuse of CCA Components!



For **Library developers**:

- + requirements explicit
- + constraints explicit
- + sequencing explicit
- + automatic enforcement
- + enhanced debugging

For **Domain Scientists**, components will be:

- + well-debugged
- + well-documented
- + easier to use

Future work focuses on adding and exploring more features.

- Add support for specifying and enforcing method sequencing
- Explore annotations and mechanisms for:
 - Enabling the use of a method within assertions (e.g., **const** or immutable)
 - Checking features of an instance (e.g., comparing sizes of two matrices or vectors)
 - Integrating and checking relevant domain-specific properties (e.g., standard units, types of matrices)
 - Automated determination of compatibility
 - Generation and automated use of translation routines
- *Anything else?*